# Projective Maximal Families of Orthogonal Measures with Large Continuum

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General Overview Some recent results Orthogonal Measures

The results which we are to consider, concern the definability of certain combinatorial objects on the real line and in particular the question of how low in the projective hierarchy such objects exist.

- (Mathias) There is no  $\Sigma_1^1$  mad family in  $[\omega]^{\omega}$ .
- (Miller) If V = L, then there is a  $\Pi_1^1$  mad family in  $[\omega]^{\omega}$ .

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# Theorem (L. Harrington)

The existence of  $\Delta_3^1$ -definable wellorder of the reals is consistent with  $\mathfrak{c}$  being as large as desired and MA.

# Theorem (S. D. Friedman)

The existence of  $\Delta_3^1$ -definable wellorder of the reals is consistent with  $\mathfrak{c} = \omega_2$  and MA.

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Note that  $\Delta_3^1$  wellorder is optimal for models of  $\mathfrak{c} > \aleph_1$ , since by a result of Mansfield if there is a  $\Sigma_2^1$  definable w.o. on  $\mathbb{R}$ , then all reals are constructible.

General Overview Some recent results Orthogonal Measures

The existence of a  $\Delta_3^1$ -definable w.o. on the reals is consistent with each of the following:

- $\begin{aligned} & \bullet \quad (\mathsf{F}., \, \mathsf{Friedman}) \ \mathfrak{d} < \mathfrak{c} = \omega_2; \ \mathfrak{b} < \mathfrak{g} = \mathfrak{c} = \omega_2; \\ & \mathfrak{b} < \mathfrak{a} = \mathfrak{s} = \mathfrak{c} = \omega_2; \end{aligned}$
- (Friedman, Zdomskyy) the existence of a Π<sub>2</sub><sup>1</sup> definable ω-mad family on [ω]<sup>ω</sup> together with b = c = ω<sub>2</sub>;
- (F., Friedman, Zdomskyy) the existence of a Π<sup>1</sup><sub>2</sub> definable ω-mad family on [ω]<sup>ω</sup> together with b = c = ω<sub>3</sub>;
- (F., Friedman, Zdomskyy) Martin's Axiom and  $\mathfrak{c} = \omega_3$ .

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General Overview Some recent results Orthogonal Measures

Let X be a Polish space, P(X) be the Polish space of Borel probability measures on X,  $\mu, \nu \in P(X)$ .

- ▶  $\mu \perp \nu$  iff  $\exists$  Borel set  $B \subseteq X$  such that  $\mu(B) = 0$  and  $\nu(X \setminus B) = 0$ .
- A set of measures A ⊆ P(X) is said to be orthogonal if whenever μ, ν ∈ A and μ ≠ ν then μ ⊥ ν.
- A maximal orthogonal family, or m.o. family, is an orthogonal family  $\mathcal{A} \subseteq P(X)$  which is maximal under inclusion.

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Theorem (Preiss, Rataj, 1985) There are no analytic m.o. families. Theorem (F., Törnquist, 2009) If V = L then there is a  $\Pi_1^1$  m.o. family.

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General Overview Some recent results Orthogonal Measures

## Theorem (F., S.D. Friedman, A. Törnquist)

There are generic extensions of L, in which there is a  $\Delta_3^1$ -definable w.o. of the reals, a  $\Pi_2^1$  definable m.o. family, there are no  $\Sigma_2^1$ -definable m.o. families and each of the following holds:

$$\bullet \ \mathfrak{d} = \omega_1 < \mathfrak{c} = \omega_2$$

 $\blacktriangleright \ \mathfrak{b} = \mathfrak{c} = \omega_3$ 

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Coding a real into a measure no  $\Sigma_2^1$ -definable m.o.families

Let X be a Polish space and let  $\mu, \nu \in P(X)$ . Then

- μ is absolutely continuous with respect to ν, written μ ≪ ν, if for all Borel B ⊆ X, if ν(B) = 0 then μ(B) = 0.
- μ, ν ∈ P(2<sup>ω</sup>) are absolutely equivalent, written μ ≈ ν, if μ ≪ ν and ν ≪ μ.

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Coding a real into a measure no  $\Sigma_2^1$ -definable m.o.families

- ▶ For  $s \in 2^{<\omega}$ , let  $N_s = \{x \in 2^{\omega} : s \subseteq x\}$  and let  $p(2^{\omega})$  be the set of all  $f : 2^{<\omega} \rightarrow [0,1]$  such that  $f(\emptyset) = 1 \land (\forall s \in 2^{<\omega})f(s) = f(s^{\frown}0) + f(s^{\frown}1).$
- p(2<sup>ω</sup>) and P(2<sup>ω</sup>) are isomorphic via f → μ<sub>f</sub> where μ<sub>f</sub> ∈ P(2<sup>ω</sup>) is uniquely determined by μ<sub>f</sub>(N<sub>s</sub>) = f(s) for all s ∈ 2<sup><ω</sup>. If μ = μ<sub>f</sub>, then f is called the code for μ.

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Coding a real into a measure no  $\Sigma_2^1$ -definable m.o.families

We describe now a way of coding  $z \in 2^{\omega}$  into a  $\mu \in P_c(2^{\omega})$ .

- ▶ For  $s \in 2^{<\omega}$ , let  $t(s, \mu)$  be the lexicographically least  $t \in 2^{<\omega}$  such that  $s \subseteq t$ ,  $\mu(N_{t\cap 0}) > 0$  and  $\mu(N_{t\cap 1}) > 0$ .
- Recursively define  $(t_n^{\mu})_{n \in \omega} \subseteq 2^{<\omega}$  by letting  $t_0^{\mu} = \emptyset$  and  $t_{n+1}^{\mu} = t(t_n^{\mu \frown} 0, \mu)$ .
- ▶ For  $f \in p_c(2^{\omega})$  and  $n \in \omega \cup \{\infty\}$  write  $t_n^f$  for  $t_n^{\mu_f}$ .

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Coding a real into a measure no  $\Sigma_2^1$ -definable m.o.families

Define the relation  $R \subseteq p_c(2^{\omega}) \times 2^{\omega}$  as follows: R(f, z) holds iff for all  $n \in \omega$  we have

$$(z(n) = 1 \longleftrightarrow (f(t_n^{f^{\frown}} 0) = \frac{2}{3}f(t_n^f) \land f(t_n^{f^{\frown}} 1) = \frac{1}{3}f(t_n^f))) \land$$
$$(z(n) = 0 \Leftrightarrow f(t_n^{f^{\frown}} 0) = \frac{1}{3}f(t_n^f) \land f(t_n^{f^{\frown}} 1) = \frac{2}{3}f(t_n^f)).$$

Whenever  $(f, z) \in R$  we say that f codes z.

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## Lemma (F., Törnquist)

There is a recursive function  $\overline{r} : p_c(2^{\omega}) \times 2^{\omega} \to p_c(2^{\omega})$  such that for all  $f \in p_c(2^{\omega})$  and  $z \in 2^{\omega}$  we have:

 $\mu_{\bar{r}(f,z)} \approx \mu_f$  and  $R(\bar{r}(f,z),z)$ .

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Coding a real into a measure no  $\Sigma_2^1$ -definable m.o.families

## Proposition

Let  $a \in \mathbb{R}$ . If there is either a Cohen or a random real over L[a], then there is no  $\Sigma_2^1(a)$  m.o. family.

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Shooting clubs Localization Coding with perfect trees Preliminaries The wellorder

We proceed with the construction of a generic extension of L in which there is a  $\Delta_3^1$  definable well order of the reals, there is a  $\Pi_2^1$ -definable m.o. family, there are no  $\Sigma_2^1$ -definable m.o. families and  $\mathfrak{d} = \omega_1 < \mathfrak{c} = \omega_2$ .

A transitive ZF<sup>-</sup> model is suitable if  $\omega_2^{\mathcal{M}}$  exists and  $\omega_2^{\mathcal{M}} = \omega_2^{L^{\mathcal{M}}}$ .

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If  $S \subseteq \omega_1$  is a stationary, co-stationary set, let Q(S) be the poset of all countable closed subsets of  $\omega_1 \setminus S$  with the extension relation end-extension.

Then Q(S) is  $\omega_1 \setminus S$ -proper,  $\omega$ -distributive and adds a club disjoint from S.

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Let  $X \subseteq \omega_1$ ,  $\phi(\omega_1, X)$  a  $\Sigma_1$ -sentence with parameters  $\omega_1, X$  which is true in all suitable models containing  $\omega_1$  and X as elements. Let  $\mathcal{L}(\phi)$  be the poset of all  $r : |r| \to 2$ , where  $|r| \in Lim(\omega_1)$  such that 1. if  $\gamma < |r|$  then  $\gamma \in X$  iff  $r(2\gamma) = 1$ , 2. if  $\gamma \leq |r|$ ,  $\mathcal{M}$  is a countable, suitable model with  $r \upharpoonright \gamma \in \mathcal{M}$ ,

$$\gamma = \omega_1^{\mathcal{M}}$$
, then  $\mathcal{M} \vDash \phi(\gamma, X \cap \gamma)$ .

The extension relation is end-extension.

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- $\mathcal{L}(\phi)$  has a countably closed dense subset,
- ▶ if G is  $\mathcal{L}(\phi)$ -generic,  $\mathcal{M}$  is countable suitable with  $(\bigcup G) \upharpoonright \gamma \in \mathcal{M}$ , where  $\gamma = \omega_1^{\mathcal{M}}$ , then  $\mathcal{M} \vDash \phi(\gamma, X \cap \gamma)$ .

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Let  $Y \subseteq \omega_1$  be generic over *L*, in *L*[*Y*] cofinalities have not been changed. Let  $\overline{\mu} = {\mu_i}_{i \in \omega_1}$  of *L*-countable ordinals such that  $\mu_i$  is

 $\min\{\mu: \mu > \sup_{j < i} \mu_j, L_{\mu}[Y \cap i] \vDash ZF^-, L_{\mu} \vDash \omega \text{ is the largest cardinal}\}.$ 

A real *R* codes *Y* below *i* if  $\forall j < i, j \in Y$  iff  $L_{\mu_i}[Y \cap j, R] \models ZF^-$ .

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- Let C(Y) be the p.o. of all perfect trees T ⊆ 2<sup><ω</sup> whose branches codes Y below |T| = min{i : T ∈ L<sub>µi</sub>[Y ∩ i]}.
- For  $T_0$ ,  $T_1$  in  $\mathcal{C}(Y)$ , let  $T_0 \leq T_1$  iff  $T_0 \subseteq T_1$ .
- C(Y) is proper and  ${}^{\omega}\omega$ -bounding.
- $\mathcal{C}(Y)$  adds a real R which codes Y.

Shooting clubs Localization Coding with perfect trees **Preliminaries** The wellorder

Fix a bookkeeping function  $F: Lim'(\omega_2) \to L_{\omega_2}$  and a sequence  $\vec{S} = (S_{\beta} : \beta < \omega_2)$  of almost disjoint stationary subsets of  $\omega_1$ , which are  $\Sigma_1$ -definable over  $L_{\omega_2}$  with parameter  $\omega_1$  and

- $F^{-1}(a)$  is unbounded in  $Lim'(\omega_2)$  for every  $a \in L_{\omega_2}$
- ▶ if  $\mathcal{M}, \mathcal{N}$  are suitable such that  $\omega_1^{\mathcal{M}} = \omega_1^{\mathcal{N}}$  then  $F^{\mathcal{M}}, \vec{S}^{\mathcal{M}}$ agree with  $F^{\mathcal{N}}, \vec{S}^{\mathcal{N}}$  on  $\omega_2^{\mathcal{M}} \cap \omega_2^{\mathcal{N}}$ .

Also if  $\mathcal{M}$  is suitable and  $\omega_1^{\mathcal{M}} = \omega_1$  then  $F^{\mathcal{M}}, \overline{S}^{\mathcal{M}}$  equal the restrictions of F,  $\vec{S}$  to the  $\omega_2$  of  $\mathcal{M}$ .

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Shooting clubs Localization Coding with perfect trees Preliminaries The wellorder

- Assume that all names for reals are nice and that for α < β < ω<sub>2</sub>, all ℙ<sub>α</sub>-names for reals precede in the canonical wellorder <<sub>L</sub> of L all ℙ<sub>β</sub>-names for reals, which are not ℙ<sub>α</sub>-names.
- ▶ If  $x \in L[G_{\alpha}] \cap {}^{\omega}\omega$  let  $\sigma_{x}^{\alpha}$  be the  $<_{L}$ -least  $\mathbb{P}_{\gamma}$ -name for x, where  $\gamma \leq \alpha$  is least so that x has a  $\mathbb{P}_{\gamma}$ -name.
- If  $x, y \in L[G_{\alpha}] \cap {}^{\omega}\omega$  let  $x <_{\alpha} y$  iff  $\sigma_{x}^{\alpha} <_{L} \sigma_{y}^{\alpha}$ .
- ▶ Then  $<^{G} = \bigcup_{\alpha < \omega_{2}} \dot{<}^{G}_{\alpha}$  will be the desired wellorder of the reals in L[G].

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For every  $\alpha \in \omega_2$ , let

- $W_{\alpha}$  be the *L*-least subset of  $\omega_1$  coding  $\alpha$ .
- $\dot{F}^0_{\alpha}$ ,  $\dot{F}^1_{\alpha}$  be  $\mathbb{P}_{\alpha}$ -names for nicely definable bijections

$$F^0_{lpha}: 2^\omega 
ightarrow p_c(2^\omega), F^1_{lpha}: (2^\omega)^\omega 
ightarrow 2^\omega$$

such that for all  $i \in \{0, 1\}$  and  $\alpha < \beta < \omega_2$  in  $L^{\mathbb{P}_\beta}$  we have  $F^i_{\alpha} \subseteq F^i_{\beta}$  (e.g. take  $(F^0_{\alpha})^{-1}$ ,  $F^1_{\alpha}$  to be Cantor diagonalization).

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We proceed with the definition of the poset. Let  $\mathbb{P}_0$  be the trivial poset. Suppose  $\mathbb{P}_{\alpha}$ ,  $\langle O_{\gamma} : \gamma < \alpha \rangle$  and  $\langle A_{\gamma} : \gamma < \alpha \rangle$  have been defined. Let  $G_{\alpha}$  be a  $\mathbb{P}_{\alpha}$ -generic filter.

#### Suppose

$$\alpha \in Lim'(\omega_2) = \{ \alpha \in Lim(\omega_2) : \alpha = \omega \cdot \omega \cdot \alpha'' \text{ for some } \alpha'' \ge 0 \}.$$

We will define  $\mathbb{P}_{\alpha+\gamma}$  for  $\gamma \in \omega \cdot \omega$  recursively as follows:

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Case A.1. If  $F(\alpha) = \{\sigma_x^{\alpha}, \sigma_y^{\alpha}\}$ , then let  $x = \sigma_x^{\alpha}[G_{\alpha}]$ ,  $y = \sigma_y^{\alpha}[G_{\alpha}]$ 

- ▶ if  $m \in \Delta(x * y)$  let  $\mathbb{Q}_{\alpha+m} = Q(S_{\alpha+m})$  and if  $m \notin \Delta(x * y)$  let  $\mathbb{Q}_{\alpha+m}$  be the random real forcing.
- ▶ In  $L^{\mathbb{P}_{\alpha+\omega}}$  let  $X_{\alpha+\omega} \subseteq \omega_1$ , coding:  $W_{\alpha}$ , (x, y), a level of L in which  $\alpha$  has size at most  $\omega_1$  and the generic  $G_{\alpha+\omega}$ , which can be regarded as a subset of an element of  $L_{\omega_2}$ .
- ► Let  $\mathbb{K}^1_{\alpha+\omega} = \mathcal{L}(\phi_{\alpha+\omega})$ , where  $\phi_{\alpha+\omega} = \phi_{\alpha+\omega}(\omega_1, X_{\alpha+\omega})$  is the  $\Sigma_1$ -sentence which holds iff  $X_{\alpha+\omega}$  codes a  $W \subseteq \omega_1$  and a pair (x, y) of reals, such that W is the *L*-least code for an ordinal  $\bar{\alpha} < \omega_2$  and  $S_{\bar{\alpha}+m}$  is non-stationary for  $m \in \Delta(x * y)$ .

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- Let Y<sub>α+ω</sub> be K<sup>1</sup><sub>α+ω</sub>-generic over L[G<sub>α+ω</sub>]. The even part of Y<sub>α+ω</sub> codes X<sub>α+ω</sub> and so codes the generic G<sub>α+ω</sub>. Then in L[Y<sub>α+ω</sub>], let K<sup>2</sup><sub>α+ω</sub> = C(Y<sub>α+ω</sub>) and let R<sub>α+ω</sub> be the real added by K<sup>2</sup><sub>α+ω</sub>.
- ► Then in particular R<sub>α+ω</sub> is added by Q<sub>α+ω</sub> = K<sup>1</sup><sub>α+ω</sub> \* K<sup>2</sup><sub>α+ω</sub> and codes the stationary kill corresponding to x < y.</p>
- For every  $\gamma \in [\alpha + \omega + 1, \alpha + \omega \cdot \omega)$  let  $\mathbb{Q}_{\alpha+\gamma}$  be the random real forcing.

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*Case A.2.* If  $F(\alpha) = \{\sigma_x^{\alpha}\}$ , then let  $x = \sigma_x^{\alpha}[G_{\alpha}]$ ,  $f = F_{\alpha}^0(x)$ . If f is not orthogonal to  $\bigcup_{\gamma < \alpha} O_{\gamma}$ , let  $\mathbb{Q}_{\alpha+\gamma}$  be the random real forcing, for all  $\gamma \in \omega \cdot \omega$ . If f is orthogonal to  $\bigcup_{\gamma < \alpha} O_{\gamma}$ , define  $\mathbb{Q}_{\alpha+\gamma}$  for  $\gamma \in \omega \cdot \omega$  recursively as follows:

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▶ Define Q<sub>α+ω</sub> just as in Case A.1, but instead of Δ(x ∗ y) use Δ(x) to determine which stationary sets will be destroyed. Let R<sub>α+ω</sub> be the generic real added by Q<sub>α+ω</sub>.

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Shooting clubs Localization Coding with perfect trees Preliminaries The wellorder

Suppose  $\mathbb{P}_{\alpha+\omega\cdot n+1}$  has been defined and  $\mathbb{Q}_{\alpha+\omega\cdot n}$  adds a real  $R_{\alpha+\omega\cdot n}$  generic over  $\mathcal{L}^{\mathbb{P}_{\alpha+\omega\cdot n}}$ . Define  $\mathbb{Q}_{\alpha+\omega\cdot n+m}$  for  $m \geq 1$  as follows:

- ▶ If  $m 1 \in \Delta(R_{\alpha+\omega \cdot n})$  let  $\mathbb{Q}_{\alpha+\omega \cdot n+m} = Q(S_{\alpha+\omega \cdot n+(m-1)})$  and if  $m - 1 \in \Delta(R_{\alpha+\omega \cdot n})$  let  $\mathbb{Q}_{\alpha+\omega \cdot n+m}$  be the random real forcing.
- ▶ In  $L[G_{\alpha+\omega\cdot n+\omega}]$  let  $X_{\alpha+\omega\cdot n+\omega} \subseteq \omega_1$  coding the sets  $W_{\alpha+\omega\cdot j}$ where  $j \leq n+1$ , the real  $R_{\alpha+\omega\cdot n}$ , a level of L in which  $\alpha+\omega\cdot n+\omega$  has size at most  $\omega_1$  and the generic  $G_{\alpha+\omega\cdot n+\omega}$ .

Shooting clubs Localization Coding with perfect trees Preliminaries The wellorder

• Let  $\mathbb{K}^1_{\alpha+\omega\cdot(n+1)} = \mathcal{L}(\phi^{n+1}_{\alpha})$ , where  $\phi^{n+1}_{\alpha}(\omega_1, X_{\alpha+\omega\cdot(n+1)})$  is the  $\Sigma_1$ -sentence which holds iff  $X_{\alpha+\omega\cdot(n+1)}$  codes a tuple  $\langle \bar{W}_j \rangle_{j \leq n+1}$  of subsets of  $\omega_1$  and a real z, such that  $\bar{W}_{n+1}$  is the *L*-least code for an ordinal  $\bar{\alpha} = \bar{\alpha}_{n+1}$ ,  $\bar{W}_j$  is the *L*-least code for the largest limit  $\bar{\alpha}_j$  strictly smaller than  $\bar{\alpha}_{j+1}$  for  $j \leq n$ , and for every  $m \in \Delta(z)$ , the set  $S_{\bar{\alpha}+m}$  is non-stationary.

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▶ Then in particular  $R_{\alpha+\omega\cdot(n+1)}$  is added by  $\mathbb{Q}_{\alpha+\omega\cdot(n+1)}$ .

Shooting clubs Localization Coding with perfect trees Preliminaries The wellorder

In 
$$L^{\mathbb{P}_{\alpha+\omega\cdot\omega}}$$
 let  $u_0^{\alpha} = x$ ,  $u_n^{\alpha} = R_{\alpha+\omega\cdot n}$  for  $n \ge 1$ . Let  $\vec{u}_{\alpha} = (u_n^{\alpha})_{n\in\omega}$   
and let

$$g_{\alpha} = \bar{r}(F^{0}_{\alpha+\omega\cdot\omega}(u^{\alpha}_{0}), F^{1}_{\alpha+\omega\cdot\omega}((u^{\alpha}_{n})_{n\geq 1}))$$

For every  $\gamma \in [\alpha, \alpha + \omega \cdot \omega)$  let  $O_{\gamma} = \{g_{\alpha}\}$ . For  $n \in \omega$ , let  $A_{\alpha+\omega\cdot n} = \alpha + \omega \cdot n + \omega \setminus \Delta(u_n^{\alpha})$  and for  $\gamma$  successor in  $[\alpha, \alpha + \omega \cdot \omega)$ , let  $A_{\gamma} = \emptyset$ .

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Case B. If  $\alpha \in Lim(\omega_2) \setminus Lim'(\omega_2)$ , or  $\alpha$  is a successor which can not be presented in the form  $\alpha' + \omega \cdot n + m$  for some  $\alpha' \in Lim'(\omega_2)$ ,  $n, m \in \omega$ , then let  $\mathbb{Q}_{\alpha}$  the random real forcing. Let  $O_{\alpha} = A_{\alpha} = \emptyset$ .

With this the recursive construction of  $\mathbb{P}_{\omega_2}$  is complete. In  $\mathcal{L}^{\mathbb{P}_{\omega_2}}$ , let  $O = \bigcup_{\alpha < \omega_2} O_{\alpha}$ ,  $F^0 = \bigcup_{\alpha \in \omega_2} F^0_{\alpha}$ ,  $F^1 = \bigcup_{\alpha \in \omega_2} F^1_{\alpha}$  and for  $\vec{z} = (z_n)_{n \in \omega} \in (2^{\omega})^{\omega}$  let  $\mathcal{R}(\vec{z}) = \bar{r}(F^0(z_0), F^1((z_n)_{n \ge 1}))$ .

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#### No accidental Stationary Kill

 $\begin{array}{l} \Pi_2^1 \mbox{ m.o. family} \\ \mbox{Proper decoding} \\ \mbox{No accidental stationary kill} \\ \mbox{Maximality} \\ \Delta_3^1 \mbox{ w.o. of } \mathbb{R} \end{array}$ 

# Lemma If G is $\mathbb{P}_{\omega_2}$ -generic and $\xi \in \bigcup_{\xi \in \omega_2} \dot{A}_{\xi}^G$ , then $S_{\xi}$ is stationary in L[G].

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No accidental Stationary Kill  $\Pi_2^1$  m.o. family Proper decoding No accidental stationary kill Maximality  $\Delta_3^1$  w.o. of  $\mathbb{R}$ 

#### Lemma

Let G be  $\mathbb{P}_{\omega_2}$ -generic and let  $g = \mathcal{R}(\vec{z}), \vec{z} = (z_n)_{n \in \omega}$ . Then  $g \in O$  if and only if for every countable suitable model  $\mathcal{M}$  such that  $g \in \mathcal{M}$ , there is  $\bar{\alpha} < \omega_2^{\mathcal{M}}$  such that for all  $n \in \omega$  the set  $S_{\alpha+\omega\cdot n+m}$  is non-stationary in  $(L[z_{n+1}])^{\mathcal{M}}$  for  $m \in \Delta(z_n)$ .

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No accidental Stationary Kill  $\Pi_2^1$  m.o. family **Proper decoding** No accidental stationary kill Maximality  $\Delta_3^1$  w.o. of  $\mathbb{R}$ 

#### Proof:

If  $g \in O$ , then  $g = g_{\alpha} = \mathcal{R}(\vec{u}_{\alpha})$  for some  $\alpha$ . Let  $\mathcal{M}$  be a countable suitable such that  $g \in \mathcal{M}$ . But then  $\vec{u}_{\alpha} \in \mathcal{M}$ , and so  $Y_{\alpha+\omega\cdot n} \cap \omega_1^{\mathcal{M}} \in \mathcal{M}$  for all n. Thus  $X_{\alpha+\omega\cdot n} \cap \omega_1^{\mathcal{M}}$  is also an element of  $\mathcal{M}$ . By definition of  $\mathcal{L}(\phi_{\alpha+\omega\cdot n}^n)$ , the set  $X_{\alpha+\omega\cdot n} \cap \omega_1^{\mathcal{M}}$  codes a tuple  $\langle W_j^n \rangle_{j \leq n}$  of subsets of  $\omega_1$  such that  $W_n^n$  is the *L*-least code of an ordinal  $\alpha_n^n$  in  $\omega_2$  and for j < n the set  $W_j^n$  is the *L*-least code for the largest limit ordinal  $\alpha_j^n$  below  $\alpha_{j+1}^n$ . It remains to observe that  $W_j^n = W_j^m$  for  $j \leq n < m$  and so  $\alpha_0^n$  does not depend on n (i.e. we have proper decoding). But then  $\bar{\alpha} = \alpha_0^n$  is the desired ordinal.

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 Introduction
 No accidental Stationary Kill

 Maximal Families of Orthogonal Measures
 Proper decoding

 The forcing construction
 No accidental stationary kill

 Maximal Families of Orthogonal Measures
 No accidental stationary kill

 Maximal Families of Orthogonal Measures
 No accidental stationary kill

 Maximal Families of Orthogonal Measures
 No accidental stationary kill

 Maximality
 ∆<sup>1</sup>/<sub>3</sub> w.o. of ℝ

Suppose that for every countable suitable model  $\mathcal{M}$  such that  $g \in \mathcal{M}$ , there is  $\bar{\alpha} < \omega_2^{\mathcal{M}}$  with the desired properties. By the Löwenheim-Skolem theorem, the same holds in  $\mathbb{H}_{\Theta}^{\mathbb{P}_{\omega_2}}$  for some large  $\Theta$ . Therefore there is  $\alpha < \omega_2^{\mathcal{M}}$  such that for all  $n \in \omega$ , the set  $S_{\alpha+\omega\cdot n+m}$  is non-stationary iff  $m \in \Delta(z_n)$ . Since there is no accidental stationary kill,  $z_n = u_n^{\alpha}$  for all n, which implies that  $g = \mathcal{R}(\vec{u}_{\alpha}) = g_{\alpha} \in O$ .  $\Box$ 

We will show that O is maximal in  $p_c(2^{\omega})$ . Suppose in  $L^{\mathbb{P}_{\omega_2}}$  there is a code f of a measure orthogonal to every measure in the family  $\overline{O} = \{\mu_g : g \in O\}$ . Choose  $\alpha$  minimal in  $Lim'(\omega_2)$  such that  $f \in L[G_{\alpha}]$  and let  $x = (F_{\alpha}^0)^{-1}(f)$ . Since  $F^{-1}(\sigma_x^{\alpha})$  is unbounded, there is  $\alpha' \geq \alpha$  in  $Lim'(\omega_2)$  such that  $F(\alpha') = \sigma_x^{\alpha}(=\sigma_x^{\alpha'})$ . But then  $g_{\alpha'}$  is a code of a measure equivalent to  $\mu_f$ , which is a contradiction. To obtain a  $\Pi_2^1$ -definable m.o. family in  $L^{\mathbb{P}_{\omega_2}}$ , consider the union of  $\overline{O}$  with the set of all point measures.

	No accidental Stationary Kill
Introduction	$\Pi_2^1$ m.o. family
Maximal Families of Orthogonal Measures	Proper decoding
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The definitions	Maximality
	$\Delta^1_3$ w.o. of $\mathbb R$

Similarly one can show that < has a  $\Delta_3^1$  definition. More precisely, we have:

#### Lemma

Let G be  $\mathbb{P}_{\omega_2}$ -generic and let x, y be reals in L[G]. Then x < y iff there is a real R such that for every countable suitable model  $\mathcal{M}$ , containing R as an element there is an  $\bar{\alpha} < \omega_2^{\mathcal{M}}$  such that  $S_{\bar{\alpha}+m}$  is non-stationary iff  $m \in \Delta(x * y)$ .

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- Since for every real  $a \in L^{\mathbb{P}_{\omega_2}}$  there is a random real over *L*, in  $L^{\mathbb{P}_{\omega_2}}$  there are no  $\Sigma_2^1$  m.o. families.
- The dominating number ∂ remains ω<sub>1</sub> in L<sup>P</sup>ω<sub>2</sub>, since the countable support iteration of S-proper <sup>ω</sup>ω-bounding posets is <sup>ω</sup>ω-bounding.