

Independence Relations

Hans Scheuermann

January 26, 2004

Organisation

- Examples (3 slides)
- Preview (1 slide)
- Axioms of Independence (5 slides)
- \wp -Forking Independence \perp^{\wp} (2 slides)
- Proof of Theorem on \perp^{\wp} (5 slides)

Examples

(1/3)

Vector space

A, B, C sets of vectors.

$\langle AB \rangle = \text{span of } A \cup B.$

$$A \underset{C}{\perp} B \quad : \Leftrightarrow \quad \langle AC \rangle \cap \langle BC \rangle = \langle C \rangle$$

$$\Leftrightarrow \begin{cases} \bar{a} \in A \text{ lin. indep. / } \langle C \rangle \\ \Rightarrow \bar{a} \text{ lin. indep. / } \langle BC \rangle \end{cases}$$

Examples

(2/3)

Algebraically closed field

A, B, C sets of field elements.

$\text{acl}(AB) =$ smallest alg. closed field $\supseteq A \cup B$.

$$A \underset{C}{\downarrow} B \quad : \Leftrightarrow \quad \begin{cases} \bar{a} \in A \text{ alg. indep. / } \text{acl}(C) \\ \Rightarrow \bar{a} \text{ alg. indep. / } \text{acl}(BC) \end{cases}$$

$$\Rightarrow \quad \text{acl}(AC) \cap \text{acl}(BC) = \text{acl}(C)$$

More complicated—not ‘modular’!

Examples

(3/3)

Everywhere infinite forest

Disjoint union of infinitely many unordered trees.

Each tree vertex has infinite degree.

A, B, C sets of vertices.

$d(a, \text{conv}(B))$ = distance between a and the convex hull of B . (May be ∞ !)

$$A \downarrow_C B \quad :\Leftrightarrow \quad \begin{cases} d(a, \text{conv}(BC)) = d(a, \text{conv } C) \\ \text{for all } a \in A \end{cases}$$

$$\Rightarrow \quad \text{conv}(AC) \cap \text{conv}(BC) = \text{conv}(C)$$

There is no such thing as a basis!

Preview

- 10 axioms for 'independence relations'.
- Definition of \perp^b .
- **Theorem**
 T complete consistent first-order theory.
 \exists independence relation for T
 $\Rightarrow \perp^b$ is an independence relation.
- Proof outline.
- Proof details.

Axioms of Independence

(1/5)

T is a complete consistent first-order theory.
 T has a big and homogeneous model.

Independence relation $A \perp_C B$:

(invariance)

$$\left. \begin{array}{l} A \perp_C B \\ (A', B', C') \equiv (A, B, C) \end{array} \right\} \Rightarrow A' \perp_{C'} B'$$

(symmetry)

$$A \perp_C B \Leftrightarrow B \perp_C A$$

...

Axioms of Independence

(2/5)

(monotonicity)

$$\left. \begin{array}{l} A \perp_C B \\ A_0 \subseteq A, B_0 \subseteq B \end{array} \right\} \Rightarrow A_0 \perp_C B_0$$

(finite character)

$$A \not\perp_C B \Rightarrow \left\{ \begin{array}{l} \exists \text{ finite } A_0 \subseteq A, B_0 \subseteq B : \\ A_0 \not\perp_C B_0 \end{array} \right.$$

(anti-reflexivity)

$$a \perp_B a \Rightarrow a \in \text{acl } B.$$

...

Axioms of Independence

(3/5)

(base monotonicity)

$$\left. \begin{array}{l} A \downarrow_C B \\ B_0 \subseteq B \end{array} \right\} \Rightarrow A \downarrow_{B_0 C} B$$

(transitivity)

$$\left. \begin{array}{l} B \downarrow_{CD} A \\ C \downarrow_D A \end{array} \right\} \Rightarrow BC \downarrow_D A$$

...

Axioms of Independence

(4/5)

(existence)

$$\forall A, B, C \quad \left\{ \begin{array}{l} \exists A' \underset{C}{\equiv} A : \\ A' \underset{C}{\perp} B. \end{array} \right.$$

(extension)

$$\left. \begin{array}{l} A \underset{C}{\perp} B \\ \hat{B} \supseteq B \end{array} \right\} \Rightarrow \left\{ \begin{array}{l} \exists A' \underset{BC}{\equiv} A : \\ A' \underset{C}{\perp} \hat{B} \end{array} \right\}$$

$$\left(\Leftrightarrow \left\{ \begin{array}{l} \exists \hat{B}' \underset{BC}{\equiv} \hat{B} : \\ A \underset{C}{\perp} \hat{B}' \end{array} \right\} \right)$$

...

Axioms of Independence

(5/5)

(local character)

$$\forall A, B \quad \exists C \subseteq B : \begin{cases} A \downarrow_C B \\ |C| < \kappa(|A|) \end{cases}$$

\mathfrak{b} -Forking Independence $\perp^{\mathfrak{b}}$ (1/2)

Definition

$$A \underset{C}{\perp^{\mathfrak{a}}} B \quad :\Leftrightarrow \quad \text{acl}(AC) \cap \text{acl}(BC) = \text{acl}(C)$$

$$A \underset{C}{\perp^{\mathfrak{m}}} B \quad :\Leftrightarrow \quad \left\{ \begin{array}{l} C \subseteq D \subseteq \text{acl}(BC) \\ \Rightarrow A \underset{D}{\perp^{\mathfrak{a}}} B \end{array} \right.$$

$$A \underset{C}{\perp^{\mathfrak{b}}} B \quad :\Leftrightarrow \quad \left\{ \begin{array}{l} \hat{B} \supseteq B \\ \Rightarrow \exists A' \underset{BC}{\equiv} A : A' \underset{C}{\perp^{\mathfrak{m}}} \hat{B} \end{array} \right.$$

Alf Onshuus:

Thorn-forking in rosy theories (Berkeley 2002)

Theorem

If there is any independence relation at all, then $\perp^{\mathfrak{b}}$ is one.

\mathfrak{p} -Forking Independence $\perp^{\mathfrak{p}}$ (2/2)

Proof Outline

1. Delete existence and symmetry axioms—
they are redundant!
2. $\perp^{\mathfrak{a}}$ satisfies all axioms except:
base monotonicity.
3. $\perp^{\mathfrak{m}}$ satisfies all undeleted axioms except:
extension, local character.
4. $\perp^{\mathfrak{b}}$ satisfies all undeleted axioms except:
local character.

S'pose there is an independence relation \perp .
Easy:

$$A \underset{C}{\perp} B \quad \Rightarrow \quad A \underset{C}{\overset{\mathfrak{b}}{\perp}} B.$$

So $\perp^{\mathfrak{b}}$ satisfies local character as well. \square

Proof of Theorem

(1/5)

1a. Existence is redundant

(existence)

$$\forall A, B, C \quad \left\{ \begin{array}{l} \exists A' \equiv_C A : \\ A' \downarrow_C B. \end{array} \right.$$

Local character $\Rightarrow A \downarrow_D C$ for some $D \subseteq C$.

Base monotonicity $\Rightarrow A \downarrow_C C$.

Extension $\Rightarrow \exists A' \equiv_C A$ such that $A' \downarrow_C BC$.

Monotonicity $\Rightarrow A' \downarrow_C B$. □

Proof of Theorem

(2/5)

1b. Symmetry is redundant

(symmetry)

$$A \downarrow_C B \quad \Rightarrow \quad B \downarrow_C A$$

Enumerate A : \bar{a}_0 .

Let $\kappa \geq \kappa(|B|)$ be regular.

Find BC -indiscernible sequence $(\bar{a}_i)_{i < \kappa}$
such that $\{\bar{a}_i \mid i < \lambda\} \downarrow_C \bar{a}_\lambda$ for $\lambda < \kappa$.

Local character

$$\Rightarrow \exists D \subseteq C\{\bar{a}_i \mid i < \kappa\} \text{ s.t. } B \downarrow_D C\{\bar{a}_i \mid i < \kappa\}.$$

κ regular and $|D| < \kappa$

$$\Rightarrow D \subseteq C\{\bar{a}_i \mid i < \lambda\} \text{ for some } \lambda < \kappa$$

$$\Rightarrow B \downarrow_{C\{\bar{a}_i \mid i < \lambda\}} C\{\bar{a}_i \mid i < \kappa\} \text{ (base monot.)}$$

$$\Rightarrow B \downarrow_{C\{\bar{a}_i \mid i < \lambda\}} \bar{a}_\lambda \quad \text{(monotonicity)}$$

$$\Rightarrow B \downarrow_C \bar{a}_\lambda \quad \text{(transitivity)}$$

$$\Rightarrow B \downarrow_C \bar{a}_0 \quad \text{(invariance) } \square$$

Proof of Theorem

(3/5)

2. Properties of \perp^a

\perp^a satisfies all axioms except:
base monotonicity.

Most axioms very straightforward.

Existence is existence of free amalgams over acl-closed sets (standard fact).

Extension follows from symmetry, existence and transitivity.

Proof of Theorem

(4/5)

3. Properties of \perp^m

\perp^m satisfies all undeleted axioms except:
extension, local character.

Base monotonicity was built into \perp^m .

Each of the other axioms is straightforward from the corresponding axiom for \perp^a .

Proof of Theorem

(5/5)

4. Properties of \perp^b

\perp^b satisfies all undeleted axioms except:
local character.

Extension was built into \perp^m .

Finite character is a bit tricky.

Each of the other axioms is straightforward from the corresponding axiom for \perp^a .