

Winter 2014, Introduction to mathematical logic

Series 2

Exercise 1 Let L be a language. Let T be a set of L -formulas which is *deductively closed* in the sense that $\varphi \in T$ whenever $T \vdash \varphi$. Let $Term_L$ be the set of all L -terms.

Show that

$$t \sim t' \iff t \doteq t' \in T$$

defines an equivalence relation on $Term_L$.

For $t \in Term_L$ let $[t] := \{t' \in Term_L \mid t \sim t'\}$ denote the equivalence class of t . Define an L -structure $\mathfrak{T}(T)$ with universe $\{[t] \mid t \in Term_L\}$ as follows:

$$\begin{aligned} ([t_1], \dots, [t_r]) \in R^{\mathfrak{T}(T)} &\iff Rt_1 \cdots t_r \in T, \\ f^{\mathfrak{T}(T)}([t_1], \dots, [t_r]) &:= [ft_1 \cdots t_r], \\ c^{\mathfrak{T}(T)} &:= [c], \end{aligned}$$

where $r \geq 1$ is a natural, $R \in L$ is an r -ary relation symbol, $f \in L$ is an r -ary function symbol and $c \in L$ is a constant $c \in L$.

Show that $\mathfrak{T}(T)$ is well-defined.

Let β be an assignment in $\mathfrak{T}(T)$ such that $\beta(x) = [x]$ for every variable x . Show that

1. $t^{\mathfrak{T}(T)}[\beta] = [t]$ for every $t \in Term_L$;
2. $\mathfrak{T}(T) \models \varphi[\beta] \iff \varphi \in T$ for every atomic L -formula φ ;

Exercise 2 This continues the above exercise. Prove:

1. If γ is an assignment in A such that $\mathfrak{A} \models T[\gamma]$, then

$$h([t]) := t^{\mathfrak{A}}[\gamma]$$

defines an homomorphism from $\mathfrak{T}(T)$ into \mathfrak{A} .

2. $\mathfrak{T}(T) \models \varphi$ if $\varphi \in T$ is an L -sentence of the form $\forall \bar{x}((\varphi_0 \wedge \dots \wedge \varphi_\ell) \rightarrow \varphi_{\ell+1})$ where each φ_i is an atomic L -formul;

3. In particular, if $L = \{-^1, e, \circ\}$ is the language of groups and T is the set of L -formulas logically implied by the group axioms, then $\mathfrak{I}(T)$ is a group such that for every finite or countable group \mathfrak{G} there is a *surjective* homomorphism from $\mathfrak{I}(T)$ onto \mathfrak{G} .

Exercise 3 Let L be a language and T an L -theory. Call an L -sentence T -*provable* if it is an element of T or a tautology or an equality axiom or an \exists -axiom or can be obtained from T -provable formulas via Modus Ponens or \exists -introduction.

Show that an L -sentence φ T -provable if and only if $T \vdash \varphi$.