## Winter 2015, Introduction to mathematical logic

Series 3

**Exercise 1** Let  $\Gamma$  be a theory. Let T be the set of terms in the language  $\mathcal{L}(\Gamma)$  of  $\Gamma$ . Then

$$t \sim t' \quad \Longleftrightarrow \quad \Gamma \vdash t = t'$$

defines an equivalence relation on T. For  $t \in T$  let

$$[t] := \{t' \in T \mid t \sim t'\}$$

denote the equivalence class of t. Define a structure  $\mathcal{T}_{\Gamma}$  for  $\mathcal{L}(\Gamma)$  with universe

$$\{[t] \mid t \in T\}$$

as follows:

$$([t_1], \dots, [t_r]) \in R^{\prime_{\Gamma}} \iff \Gamma \vdash Rt_1 \cdots t_r,$$
  
$$f^{\mathcal{T}_{\Gamma}}([t_1], \dots, [t_r]) := [ft_1 \cdots t_r],$$
  
$$c^{\mathcal{T}_{\Gamma}} := [c],$$

where  $r \in \mathbb{N}$  is a natural, R is an r-ary relation symbol of  $\mathcal{L}(\Gamma)$ , f is an r-ary function symbol of  $\mathcal{L}(\Gamma)$  and c is a constant of  $\mathcal{L}(\Gamma)$ .

Prove:

- (a)  $\mathcal{T}_{\Gamma}$  is well-defined;
- (b) if  $t \in T$  has variables  $\{x_1, \ldots, x_k\}$  and  $t_1, \ldots, t_k \in T$ , then

$$t(x_1/[t_1],\ldots,x_k/[t_k])^{\mathcal{T}_{\Gamma}} = [t(x_1/t_1,\ldots,x_k/t_k)];$$

(c) if  $\varphi$  is an atomic  $\mathcal{L}(\Gamma)$ -formula with free variables  $\{x_1, \ldots, x_k\}$  and  $t_1, \ldots, t_k \in T$ , then

$$\mathcal{T}_{\Gamma} \models \varphi(x_1/[t_1], \ldots, x_k/[t_k]) \Longleftrightarrow \Gamma \vdash \varphi(x_1/t_1, \ldots, x_k/t_k).$$

**Exercise 2** This continues the above exercise. Prove:

(a) Let  $\mathcal{A} \models \Gamma$ . Let  $\gamma$  be a map from the set of variables into  $\mathcal{A}$ . Set

$$h([t]) := t(x_1/\gamma(x_1), \dots, x_k/\gamma(x_k))^{\mathcal{A}}$$

for each  $t \in T$  with free variables  $\{x_1, \ldots, x_k\}$ . Show that h is a homomorphism from  $\mathcal{T}_{\Gamma}$  into  $\mathcal{A}$ .

(b) Let  $\varphi$  be an  $\mathcal{L}(\Gamma)$ -sentence of the form  $\forall \bar{x}((\varphi_0 \land \ldots \land \varphi_\ell) \to \varphi_{\ell+1})$  where each  $\varphi_i$  is atomic. Show:

if 
$$\Gamma \vdash \varphi$$
, then  $\mathcal{T}_{\Gamma} \models \varphi$ .

(c) Let  $\mathcal{L}$  be the first-order language containing function symbols + (binary), - (unary) and a constant 0. Give a theory  $\Gamma$  in this language whose models are precisely the groups (with + denoting the group operation, - the inverse and 0 the neutral element).

Show that  $\mathcal{T}_{\Gamma}$  is a group such that for every finite or countable group  $\mathcal{G}$  there is a surjective homomorphism from  $\mathcal{T}_{\Gamma}$  onto  $\mathcal{G}$ .

**Exercise 3** Let  $\mathcal{L}$  be as in the previous exercise (c). Show:

- (a) There exists a theory  $\Gamma'$  whose models are precisely the infinite groups.
- (b) If  $\Gamma''$  is a theory whose models contain arbitrarily large finite groups, then there is an infinite group that is a model of  $\Gamma''$ .
- (c) An  $\mathcal{L}$ -sentence that holds in all infinite groups, also holds in some finite group. Conclude that a theory  $\Gamma'$  as in (a) cannot be finite.

**Exercise 4** Let  $\mathcal{L}$  be an arbitrary first-order language. Let  $\Gamma, \Gamma'$  be theories in this language. Assume  $\Gamma \cup \Gamma'$  does not have a model. Show that there exists an  $\mathcal{L}$ -sentence  $\varphi$  such that  $\Gamma \vdash \varphi$  and  $\Gamma' \vdash \neg \varphi$ .